Parallel Reduction with CUDA

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Input: Array of ints a[0], a[1], ..., a[n-1]

1 4 3 2 8 6 3 2 1 0 3 2 1 3 2 3 2 9 1 2 3 4 5 6 1 1 2 3 0 0 2 1

Output: The sum of the elements

```
int sum = 0;
for( int i = 0 ; i < N ; ++i )
  sum += a[i];</pre>
```

Outline of the Algorithm



1st Step: Each thread computes its partial sum

2nd Step: Each block of threads computes its partial sum

3rd Step: One block computes the final sum

Blocks and Threads



In those slides we use 2 blocks and 4 threads per block





1ST STEP



Each thread initializes its partial sum to 0



int my_sum = 0;

Each thread reads its element from a and update its sum



1st Step: Threads Compute Their Partial Sums



Each thread moves to its next element

And updates its partial sum





And so on, until all threads have their partial sums



1st Step: Threads Compute Their Partial Sums



Each thread writes its partial sum to shared memory



- Shared memory is shared by threads in the same block
- No synchronization: Threads do not know what others are doing

Advanced note: Inside a block, threads of a same warp are implicitly synchronized. On current HW a warp contains 32 threads.



We want smem to contain all the sums of the threads in the block

syncthreads() to create a synchronization barrier

Block 0 Block 0 Block 0 Block 0 13 11 13 11 13 11 11 9 smem: 0, 2: complete their writes 0, 2: wait at the barrier 0, 2, 3: wait at the barrier Threads go through the barrier 3: completes its write 4: completes its write Time

We use

Note: ____syncthreads synchronizes the threads of a same block. There is no instruction to synchronize threads of different blocks from a kernel.

1st Step: Threads Compute Their Partial Sums



Code summary









With more threads and more blocks



Half of the threads updates the sum in smem





Time

We want smem to contain all the sums: _____syncthreads()

Time

1/4th of the threads updates the sum in smem

- We keep doing that until there is no more work to do
- Code summary

Block leaders stores the sum of the block in global memory

1st and 2nd steps are implemented in one CUDA kernel

We launch the kernel

3rd Step: 1st Block Computes The Sum

We can call the same kernel with different arguments

The result is in partial sums[0]

Parallel Scan with CUDA

Input: Array of ints a[0], a[1], ..., a[n-1]

1 4 3 2 8 6 3 2 1 0 3 2 1 3 2 3 2 9 1 2 3 4 5 6 1 1 2 3 0 0 2 1

Output: The sums b[i] = a[0] + ... + a[i-1]

```
int sum = 0;
for( int i = 0 ; i < N ; ++i )
{
    b[i] = sum;
    sum += a[i];
}</pre>
```

Outline of the Algorithm

- 1st Step: Each block computes its partial sum
 - Same kernel as the 1st kernel of the reduction
- 2nd Step: One block scans the block sums (global scans)
- 3rd Step: Each block scans its elements and uses its global scan

Scan Primitive: Parallel Scan in a Block

Each thread loads its item from GMEM

int my_sum = in_buffer[idx];

we allocate a	buffer in SMEN	of size 2 x	blockDim.x

	0	1	2	3	4	5	6	7
smem:	0	0	0	0	2	3	1	4

my sum:

2

3

Store 0s in the 1st half of smem and its items in the 2nd half

Scan Primitive: Parallel Scan in a Block

Parallel scan pattern

Homework

How can you implement inclusive scan?

```
int sum = 0;
for( int i = 0 ; i < N ; ++i )
{
    b[i] = sum += a[i];
}</pre>
```

How can you optimize your implementation?

Parallel Radix Sort with CUDA

Outline of the Algorithm

- Loop on the bits:
- 1st Step: Each block computes its partial sums
 - One sum for each combination of bits
 - Almost the same kernel as the 1st kernel of the reduction
- 2nd Step: One block scans the block sums (global scans)
- 3rd Step: Each block scatters its elements

Sort Primitive: Where to Store Counters

Each thread has one counter per bit combination

int my_counters[NUM_COUNTERS];

Problem: It's not possible to dynamically address registers

my_counters[(item & mask) >> i]++;

Solution: Use SMEM to store counters

shared int s_counters[NUM_COUNTERS][BLOCK_DIM];

s counters[(item & mask) >> i][threadIdx.x]++;